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On a Characterization of Bilinear Forms Graphs

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We show that the bilinear forms graphs $H_q(n, d)$ of diameter $d \geq 3$ are characterized as distance-regular graphs by their parameters provided that either $n \geq d + 3$ and $q \geq 3$, or $n \geq d + 4$ and $q = 2$. As a corollary of the method used, we can show the following. If Γ is a distance-regular graph with classical parameters (d, q, α, β) and diameter $d \geq 3$, then either Γ is a Johnson graph, a Grassmann graph, a Hamming graph, or a bilinear forms graph, or β is bounded in terms of d, q and α .

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1. INTRODUCTION

We consider only finite, simple and undirected graphs Γ . The *distance* of two vertices v and w is denoted by $d(v, w)$, and the maximum distance that occurs between two vertices of Γ is the *diameter* of Γ , denoted by d . We put $\Gamma_i(v) = \{w \mid d(v, w) = i\}$. The graph Γ is *distance regular* with *intersection array* $\{b_0, b_1, \dots, b_{d-1}; c_1, c_2, \dots, c_d\}$, if for any two vertices v and w at distance i , we have $c_i = |\Gamma_{i-1}(v) \cap \Gamma_1(w)|$, $b_i = |\Gamma_{i+1}(v) \cap \Gamma_1(w)|$, and $a_i = |\Gamma_i(v) \cap \Gamma_1(w)|$; here $c_0 := b_d := 0$ and $a_i := b_0 - b_i - c_i$. The distance-regular graph Γ has *classical parameters* (d, q, α, β) , if it has diameter d , and if

$$b_i = \left(\begin{bmatrix} d \\ 1 \end{bmatrix} - \begin{bmatrix} i \\ 1 \end{bmatrix} \right) \left(\beta - \alpha \begin{bmatrix} i \\ 1 \end{bmatrix} \right) \quad \text{and} \quad c_i = \begin{bmatrix} i \\ 1 \end{bmatrix} \left(1 + \alpha \begin{bmatrix} i-1 \\ 1 \end{bmatrix} \right)$$

for $i = 0, \dots, d$. Here $\begin{bmatrix} i \\ j \end{bmatrix}$ denotes the Gaussian coefficient with basis q ; in particular, we have $\begin{bmatrix} i \\ 1 \end{bmatrix} = \sum_{k=0}^{i-1} q^k$. Throughout, we use the abbreviation $r := \begin{bmatrix} d \\ 1 \end{bmatrix}$.

Many distance-regular graphs related to classical groups and groups of Lie type have classical parameters. A typical example is the *Grassmann graph* $J_q(n, d)$, whose vertices are the subspaces of rank d of a vector space V of rank $n \geq 2d \geq 4$ over the field F_q , where two vertices are adjacent if they meet in a subspace of V of rank $d - 1$. It has classical parameters $(d, q, q, \begin{bmatrix} n-d+1 \\ 1 \end{bmatrix} - 1)$. For integers $n \geq d \geq 2$, a vector space of rank $d + n$ and a fixed subspace N of rank n , the *bilinear forms graph* $H_q(n, d)$ is the subgraph of $J_q(n + d, d)$ induced on the vertices that meet N trivially; it has classical parameters $(d, q, q - 1, q^n - 1)$. We remark that the Johnson graph $J(n, d)$ has classical parameters $(d, 1, 1, n - d)$ ($n \geq 2d \geq 4$) and the Hamming graph $H(d, n)$ has classical parameters $(d, 1, 0, n - 1)$ ($n, d \geq 2$).

The Hamming graphs $H(d, n)$ have been characterized (by their parameters) for $d \geq 2$ and $n \neq 4$ by Egawa [5]. The Johnson graphs $J(n, d)$ have been characterized for $(d, n) \neq (2, 8)$ independently by Neumaier [10] and Terwilliger [11]. The Grassmann graphs are characterized for $n \neq 2d, 2d + 1$ and $d \geq 3$ (with some exceptions if $q = 2$ or $q = 3$) in [9]. The bilinear forms graphs have been characterized for $n \geq 2d \geq 6$ and $q \geq 4$ by Huang [7] and Cuypers [4]. Thus all these graphs have satisfactory characterizations, except that the bound $n \geq 2d$ for the bilinear forms graphs is quite weak. It is the purpose of this paper to improve this bound.

THEOREM 1.1. *Let Γ be a distance-regular graph with classical parameters (d, q, α, β) where $\alpha = q - 1$ and $d \geq 3$. Suppose that either $q = 2$ and $\beta \geq q^{d+4} - 1$ or $q \geq 3$ and $\beta \geq q^{d+3} - 1$. Then q is a prime power, $\beta = q^n$ for some integer n , and Γ is the bilinear forms graph $H_q(n, d)$.*

COROLLARY 1.2. *The bilinear forms graphs $H_q(n, d)$ of diameter $d \geq 3$ are uniquely determined as distance-regular graphs by their parameters provided that either $q = 2$ and $n \geq d + 4$ or $q \geq 3$ and $n \geq d + 3$.*

The four families of distance-regular graphs mentioned above have the property that β can be arbitrary large with respect to the other three parameters d , q , and α . The following result shows that there are no other graphs with this property.

COROLLARY 1.3. *Suppose that the distance-regular graph Γ has classical parameters (d, q, α, β) with $d \geq 3$ and that Γ is not a Grassmann graph, a bilinear forms graph, a Hamming graph, or a Johnson graph. Then β is bounded in terms of d , q and α . More precisely, if $r = 1 + q + q^2 + \dots + q^{d-1}$, then we have:*

- (a) *If $q < 1$, then $|\beta| \leq 1 + |\alpha|(|r| + |q| - 1)$.*
- (b) *If $q \geq 1$ and $\alpha \neq q, q - 1$, then $|\beta - \alpha + q| \leq r(r - 1)(q + 1)|\alpha - q| \cdot |\alpha + 1 - q|$.*
- (c) *If $q \geq 1$ and $\alpha = q$, then $\beta < 8r(q^2 + 2q)/3$.*
- (d) *If $q \geq 1$ and $\alpha = q - 1$, then $\beta < (2q^4 + 2q^3 + 2q^2 + q - 1)r/(2q - 1)$.*

REMARK 1.4. Suppose that Γ is a distance-regular graph of diameter $d = 2$, that is, Γ is a strongly regular graph. It is easy to see that Γ has classical parameters $(2, q, \alpha, \beta)$ and these can be uniquely chosen in such a way that $\beta > 0$. If the parameters are integers and if β is sufficiently large with respect to α and q , then Bose [1] (see also [8]) showed that the graph is the collinearity graph of a partial geometry with parameters $(q + 1, \beta + 1, \alpha + 1)$. This is the analogue to Corollary 1.3 in the case $d = 2$.

2. BILINEAR FORMS GRAPHS

We need the following special case of Corollary 1.3 in [8], which was proved by an improved Bose–Lasker argument. Recall that a *clique* of a graph Γ is a set of mutually adjacent vertices of Γ .

RESULT 2.1. Let $c_2 \geq 1$ and k, λ, s be integers. Suppose that Γ is a regular graph of valency k such that non-adjacent vertices have at most c_2 common neighbours, and adjacent vertices have λ common neighbours. Suppose furthermore that

- (a) $\lambda > (2s - 1)(c_2 - 1) - 1$ and
- (b) $k < (s + 1)(\lambda + 1) - \frac{1}{2}s(s + 1)(c_2 - 1)$.

Define a line to be a maximal clique C satisfying $|C| \geq \lambda + 2 - (s - 1)(c_2 - 1)$. Then every vertex is in at most s lines, and any two adjacent vertices occur together in a unique line.

For a distance-regular graph with classical parameters (d, q, α, β) with $d \geq 3$, we put

$$r := \begin{bmatrix} d \\ 1 \end{bmatrix} \quad \text{and} \quad \lambda := b_0 - a_1 - c_1 = \beta - 1 + (r - 1)\alpha.$$

Notice that every vertex has $b_0 = r\beta$ neighbours and adjacent vertices have λ common neighbours. Proposition 6.2.1 in [2] shows that q is an integer with $q \neq 0, -1$. Hence, if $\alpha = q - 1$, then

$$q \text{ is an integer, } q \neq 0, -1 \text{ and } c_2 = q(q + 1) \geq 2.$$

PROPOSITION 2.2. *Let Γ be a distance-regular graph with classical parameters (d, q, α, β) of diameter $d \geq 3$ with $\alpha = q - 1 \geq 1$. Suppose that there exists an integer $s \geq r$ such that the following properties are fulfilled.*

- (a) *If $q = 2$ and $d = 3$, then $s = r = 7$.*
- (b) *$(s + 1)(\lambda + 1) - \frac{1}{2}s(s + 1)(q^2 + q - 1) > r\beta$*
- (c) *$\lambda + 1 > s(q^3 + q^2 + 2q - 1) - q^2(q^2 + q + 1)$.*

Then q is a prime power, $\beta = q^n - 1$ for some integer $n \geq d$ and Γ is the bilinear forms graph $H_q(n, d)$.

We prove this proposition in several lemmas. By a *line* we shall mean a maximal clique C of Γ satisfying $|C| \geq \lambda + 2 - (s - 1)(c_2 - 1)$.

LEMMA 2.3. *$s \geq r \geq q^2 + q + 1$ and $\lambda + 1 > s(q^3 + 2q) - q^2 - q - 1$ and $\beta > s(q^3 + q)$.*

PROOF. As $d \geq 3$, we have $s \geq r \geq q^2 + q + 1$. The bound for $\lambda + 1$ follows immediately from hypothesis (C). As $\lambda + 1 - \beta = (r - 1)(q - 1) \leq s(q - 1)$, it follows that $\beta > s(q^3 + q) + s - q^2 - q - 1 \geq s(q^3 + q)$. \square

LEMMA 2.4.

- (a) *Every vertex is on at most s lines.*
- (b) *Every two adjacent vertices v_1 and v_2 lie in a unique line, denoted by $v_1 v_2$.*
- (c) *Every line has at most $\beta + 1$ vertices.*
- (d) *Every vertex v is on at least r lines with equality iff all lines on v have $\beta + 1$ vertices.*

PROOF. As $c_2 = q^2 + q$, it follows from hypothesis (b) in Proposition 2.2 and Lemma 2.3 that conditions (a) and (b) of Result 2.1 are satisfied. This proves (a) and (b).

The eigenvalues of Γ are given by $\Theta_i := b_i/q^i - (q^i - 1)/(q - 1)$, $i = 0, \dots, d$, (see Theorem 8.4.2 in [2]). Hence, $\Theta_d = -r$ is the smallest eigenvalue of Γ . It follows from a result of Hoffman [6] (see also Proposition 4.4.6 in [2]) that every clique (and thus every line) of Γ has at most $1 - b_0/\Theta_d = \beta + 1$ vertices. This proves (c). As every vertex has $b_0 = r\beta$ neighbours, part (d) is a consequence of parts (b) and (c). \square

The *distance* from a vertex v to a line L is the minimum distance from v to a vertex of L . It will be denoted by $d(v, L)$. For two vertices v and w at distance two, we denote by $[v, w]$ the set consisting of all lines L on v with the property that $d(w, x) \leq 2$ for all $x \in L$.

LEMMA 2.5.

- (a) *If $d(v, w) = 2$, then $|[v, w]| \leq q + 1$.*
- (b) *If a vertex has distance one from a line, then it has at most $q + 1$ neighbours on that line.*
- (c) *If a vertex v lies on t lines, then $|L| \geq \lambda + 2 - (t - 1)q$ for every line L on v .*
- (d) *$|L| \geq \lambda + 2 - (s - 1)q$ for every line L .*

PROOF. By n we denote the maximum integer such that there exist vertices v_1 and v_2 at distance two with $|[v_1, v_2]| = n + 1$. As vertices at distance two have c_2 common neighbours, we have $n \leq c_2 - 1$.

Consider a vertex v_1 and a line L with $d(v_1, L) = 1$. As L is a maximal clique, there exists $v_2 \in L$ with $d(v_1, v_2) = 2$. As the lines on v_1 that meet L are in $[v_1, v_2]$, it follows that

$|\Gamma_1(v_1) \cap L| \leq |[v_1, v_2]| \leq n + 1$. Hence, every vertex that has distance one from a line has at most $n + 1$ neighbours on that line.

Consider a vertex v and suppose that v is on t lines. Consider any line L on v and a second vertex $w \in L$. Then every other line on v contains at most n neighbours of w other than v . As v lies on t lines, it follows that v and w have at most $|L| - 2 + (t - 1)n$ common neighbours. Hence $|L| \geq \lambda + 2 - (t - 1)n$ for every line L on a vertex that lies on t lines. As every vertex is on at most s lines, it follows that every line has at least $\lambda + 2 - (s - 1)n$ vertices.

It remains to show $n \leq q$. Assume on the contrary that $n \geq q + 1$. Consider vertices v and w at distance two such that $[v, w] = n + 1$, and denote by M the set consisting of the neighbours of v that have distance one or two from w . Then $|M| = b_0 - b_2 = (q + 1)\lambda - (q + 1)(q^2 - q - 1)$.

As every line has at least $\lambda + 2 - (s - 1)n$ vertices, every line of $[v, w]$ has at least $\lambda + 1 - (s - 1)n$ vertices in M . If we define $f(x) := (x + 1)(\lambda + 1 - (s - 1)x)$, then it follows that the $n + 1$ lines of $[v, w]$ cover at least $f(n)$ vertices of M . Hence $f(n) \leq |M|$. As $q + 1 \leq n \leq c_2 - 1$, we have $f(n) \geq \min\{f(q + 1), f(c_2 - 1)\}$. As $c_2 = q^2 + q$, we have $f(c_2 - 1) - f(q + 1) = (q^2 - 2)[\lambda + 1 - (q + 1)^2(s - 1)] \geq 0$. Hence $f(q + 1) \leq f(n)$ and thus

$$(q + 2)(\lambda + 1 - (s - 1)(q + 1)) = f(q + 1) \leq |M| = (q + 1)(\lambda + 1) - (q + 1)(q^2 - q - 1).$$

This is equivalent to

$$\lambda + 1 \leq s(q + 1)(q + 2) - (q + 1)(q^2 + 1).$$

Using Lemma 2.3, it follows that $q^3 + 2q < (q + 1)(q + 2)$. As $q \geq 2$, this is a contradiction. This proves that $n \leq q$. \square

LEMMA 2.6. *Every vertex is on at most $9r/8$ lines.*

PROOF. Suppose that v is on t lines. Since v has $r\beta$ neighbours, it follows from Lemma 2.5 (c) that

$$t[\lambda + 1 - (t - 1)q] \leq r\beta.$$

As $\lambda + 1 = \beta + (r - 1)(q - 1)$, this is equivalent to

$$(t - r)[\beta - t(q - 1)] \leq t(t - 1).$$

As $t \leq s$, we obtain $\beta > t(q^3 + q)$ from Lemma 2.3. It follows that $(t - r)t(q^3 + 1) \leq t(t - 1)$. As $q \geq 2$, this implies that $9(t - r) \leq t$, or $8t \leq 9r$. \square

LEMMA 2.7. *If $d(v_1, v_2) = 3$, then $|\Gamma_2(v_1) \cap L| \leq q(q^2 + q + 1)$ for every line L on v_2 .*

PROOF. Put $X = \Gamma_2(v_1) \cap L$ and $W = \Gamma_1(v_1) \cap \Gamma_2(v_2)$. If $x \in X$, then $\Gamma_1(x) \cap \Gamma_1(v_1) \subseteq W$. Hence, every vertex of X has c_2 neighbours in W . Counting pairs of adjacent vertices $x \in X$ and $w \in W$ gives

$$|X|c_2 = \sum_{w \in W} |\Gamma_1(w) \cap L|.$$

Using Lemma 2.5 (b), we obtain $|X|c_2 \leq |W|(q + 1)$. As $|W| = c_3 = q^2(q^2 + q + 1)$ and $c_2 = q^2 + q$, the assertion follows. \square

LEMMA 2.8. *If $d(v, w) = 2$, then $[v, w] = q + 1$.*

PROOF. Denote by M the set consisting of the neighbours of v that have distance one or two from w . Then $|M| = b_0 - b_2 = (q+1)(\lambda+1) - (q+1)(q^2 - q - 1)$. Put $c := q(q^2 + q + 1)$. Lemma 2.5 shows that $||[v, w]| \leq q + 1$.

Assume that $||[v, w]| \leq q$. By Lemma 2.4, every line on v contains at most $\beta + 1$ vertices and thus at most β vertices in M . A line on v that is not in $[v, w]$ contains a vertex at distance three from w and contains therefore, by Lemma 2.7, at most $c - 1$ vertices in M . As v is on at most s lines, it follows that

$$|M| \leq q\beta + (s - q)(c - 1) = q(\lambda + 1 - (r - 1)(q - 1)) + (s - q)(c - 1).$$

As $|M| = (q+1)(\lambda+1) - (q+1)(q^2 - q - 1)$, this gives

$$\lambda + 1 \leq (q+1)(q^2 - q - 1) - (r - 1)q(q - 1) + (s - q)(c - 1).$$

As $d \geq 3$, we have $r - 1 \geq q^2 + q \geq q + 1$. Hence $\lambda + 1 \leq (s - q)(c - 1)$. This contradicts hypothesis (c) in Proposition 2.2. \square

LEMMA 2.9. *Suppose that every point is on exactly r lines. Then $|\Gamma_1(v) \cap L| = q$ for every vertex v and every line L with $d(v, L) = 1$.*

PROOF. This can be shown using the argument of the proof of Lemma 2.3 in [7]. We give a slightly different proof.

As every vertex is on r lines, Lemma 2.4 shows that every line has $\beta + 1$ vertices. Fix a vertex v and denote by T the set consisting of all triples (u, u', w) of distinct pairwise collinear vertices with $d(v, u) = d(v, u') = 1$ and $d(v, w) = 2$. Furthermore, put $k_i = |\Gamma_i(v)|$ for $i = 1, 2$.

If $d(v, w) = 2$, then v and w have $q^2 + q$ common neighbours, which are on the $q + 1$ lines of $[w, v]$. It follows that w occurs in at least $(q + 1)q(q - 1)$ triples of T with equality iff every line of $[w, v]$ contains exactly q neighbours of x . Hence $|T| \geq k_2q(q^2 - 1)$.

Now consider a neighbour u of v . Then u and v have $\lambda - (\beta - 1) = (r - 1)(q - 1)$ common neighbours that do not lie on the line uv . These lie on one of the $r - 1$ lines L_2, \dots, L_r other than uv on u . If $|\Gamma_1(v) \cap L_i| = 1 + x_i$, then $\sum x_i = (r - 1)(q - 1)$ and u occurs in $\sum x_i(\beta - x_i)$ triples of T . As $\sum x_i^2 \geq (\sum x_i)^2 / (r - 1) = (r - 1)(q - 1)^2$, it follows that u occurs in at most $\beta(r - 1)(q - 1) - (r - 1)(q - 1)^2$ triples of T with equality iff all lines other than uv on u contain exactly q neighbours of v . Hence $|T| \leq k_1(r - 1)(q - 1)(\beta + 1 - q)$.

Comparing both bounds for T using $k_2 = k_1b_1/c_2$, we obtain equality. Hence, every line at distance one from v contains exactly q neighbours of v . \square

LEMMA 2.10. *Consider a vertex p and a line L with $d(p, L) = 1$. Then $[p, v_1] = [p, v_2]$ for all vertices $v_1, v_2 \in L \cap \Gamma_2(p)$.*

PROOF. Assume that there exist vertices $v_1, v_2 \in L \cap \Gamma_2(p)$ with $[p, v_1] \neq [p, v_2]$. Put $L_i := \Gamma_i(p) \cap L$ and $n_i := |L_i|$, $i = 1, 2$, and $T := [p, v_1] \cap [p, v_2]$. Then $n_1 \leq t := |T| \leq q$. Denote by T_2 the set consisting of the lines on p that are not in T . Then $|T_2| \leq s - t$, since p lies on at most s lines.

Consider a line $H \in T_2$. W.l.o.g. $H \notin [p, v_1]$. Then there exists $h \in H$ with $d(v_1, h) = 3$. Lemma 2.7 shows that $|\Gamma_2(h) \cap L| \leq c := q(q^2 + q + 1)$. As $L_1 \subseteq \Gamma_2(h) \cap L$, we obtain $|\Gamma_2(h) \cap L_2| \leq c - n_1$. It follows that there exist at most $c - n_1$ vertices $v \in L_2$ with $H \in [p, v]$. On the other hand, if $v \in L_2$, then $[p, v]$ contains at least $q + 1 - t$ lines of T_2 . Counting pairs (v, H) with $v \in L_2$ and $H \in T_2 \cap [p, v]$, we obtain

$$n_2(q + 1 - t) \leq |T_2|(c - n_1) \leq (s - t)(c - n_1).$$

As $t \leq q$ and $s \geq r > q + 1$, we have $s - t \leq (q + 1 - t)(s - q)$. It follows that $n_2 \leq (s - q)(c - n_1)$. As $d(p, L) = 1$, we have $n_1 \geq 1$. Hence $|L| = n_1 + n_2 \leq 1 + (s - q)(c - 1)$. From Lemma 2.5 we obtain $|L| \geq \lambda + 2 - (s - 1)q$. Hence

$$\lambda + 1 \leq (s - q)(c - 1) + (s - 1)q = s(c + q - 1) - qc.$$

As $c = q(q^2 + q + 1)$, this contradicts hypothesis (c) in Proposition 2.2. \square

Next we verify the dual of Pasch's Axiom. This means that for all adjacent points v and v' we have: if X is the set consisting of all vertices that are not on the line vv' and which are adjacent to v and v' , then X is a clique.

LEMMA 2.11. *The incidence structure consisting of the vertices and lines satisfies the dual of Pasch's Axiom.*

PROOF. Consider two adjacent vertices v and v' , put $H = vv'$, and denote by X the set consisting of all common neighbours of v and v' that are not on H . We have to show that X is a clique. As $|H| \leq \beta + 1$ (Lemma 2.4) and $|H| - 2 + |X| = \lambda$, we have $|X| \geq \lambda + 1 - \beta = (r - 1)(q - 1) = q^d - q$.

Consider $x \in X$. We shall first show that $|\Gamma_1(x) \cap X| \geq |X| - 1 - q(q - 1)$. Put $L = vx$. Then there exists a vertex $p \in L$ with $d(p, v') = 2$. Consider any vertex $x' \in X$ with $x' \neq x$ and put $L' := v'x'$. Then $L' \neq H$. If $L' = v'x$, then clearly $x' \sim x$. If x' is on one of the $q - 1$ lines of $[v', p]$ other than H and $v'x$, then x' might not be adjacent to x . As each of these $q - 1$ lines contains at most q common neighbours of v and v' (Lemma 2.5), this gives rise to at most $(q - 1)q$ vertices in X that are not adjacent to x . Finally, consider the case that $L' \notin [v', p]$. Then $d(x', p) \geq 2$. As $d(x', L) = 1$, it follows that $d(x', p) = 2$. This implies that $x \sim x'$ (otherwise $d(x', x) = 2$ and Lemma 2.10 would give $[x', p] = [x', x]$; however $L' \in [x', x]$, as $d(x, L') = 1$, but $L' \notin [x', p]$, as $L' \notin [v', p]$). This proves that $|\Gamma_1(x) \cap X| \geq |X| - 1 - q(q - 1)$.

Now consider two distinct vertices $x_1, x_2 \in X$, put $X_i := X \cap \Gamma_1(x_i)$, $i = 1, 2$, and assume that $x_1 \not\sim x_2$. Then $|X_i| \geq |L| - q(q - 1) - 1$ and $x_1, x_2 \notin X_1 \cup X_2$. Hence

$$\begin{aligned} |X_1 \cap X_2| &\geq |X_1| + |X_2| - |X_1 \cup X_2| \geq 2(|X| - q(q - 1) - 1) - (|X| - 2) \\ &= |X| - 2q(q - 1). \end{aligned}$$

As $v, v' \in \Gamma_1(x_1) \cap \Gamma_1(x_2)$, it follows that $|\Gamma_1(x_1) \cap \Gamma_1(x_2)| \geq |X| + 2 - 2q(q - 1)$. On the other hand x_1 and x_2 have at most c_2 common neighbours, because $d(x_1, x_2) = 2$. Hence $|X| \leq c_2 + 2q(q - 1) - 2 = 3q^2 - q - 2$. From $|X| \geq q^d - q$, we obtain $q^d \leq 3q^2 - 2$. As $d \geq 3$ and $q \geq 2$, it follows that $d = 3$ and $q = 2$.

Hypothesis (a) in Proposition 2.2 says that $s = r = 7$ in this case, that is, every vertex lies on precisely r lines. Lemma 2.9 shows that a vertex that has distance one from a line has precisely $q = 2$ neighbours on this line. For $x \in X$, we can now improve the above bound to $|\Gamma_1(x) \cap X| \geq |X| - 1 - (q - 1)^2 = |X| - 2$. As x_1 and x_2 are not adjacent, it follows that x_1 and x_2 are adjacent to every vertex of $X \setminus \{x_1, x_2\}$. Consider the line $L_1 := vx_1$. Lemma 2.9 shows that x_1 and v are the only neighbours of v' on L , and that x_2 has two neighbours on L . Hence, there exists a vertex $w \neq v$ on L that is adjacent to x_2 but not to v' . But now the vertices v, v', w and the vertices of $X \setminus \{x_1, x_2\}$ are adjacent to x_1 and x_2 . Since $c_2 = (q + 1)q = 6$ and $|X| \geq q^d - q = 6$, this is a contradiction. \square

Every maximal clique that is not a line will be called an *assembly*.

LEMMA 2.12.

- (a) Two adjacent vertices v and v' lie in a unique assembly. It contains all common neighbours of v and v' that are not on the line vv' .
- (b) If a line L and an assembly B share two vertices, then $|L| + |B| - |L \cap B| = \lambda + 2$.

PROOF. (a) Denote by X the set consisting of all common neighbours of v and v' that are not on $L := vv'$. As $\lambda = \beta - 1 + (r - 1)(q - 1) > \beta - 1 \geq |L| - 2$, we have $X \neq \emptyset$.

By the dual of Pasch's Axiom, X is a clique. The dual of Pasch's Axiom also shows that every vertex $l \in L$ that is adjacent to some vertex x of X is also adjacent to every other vertex y of X , because l and y do not lie on one of the lines xv and xv' . This shows that $X \cup \{v, v'\}$ lies in a unique maximal clique B , which consists of X and all vertices of L that are adjacent to one (and thus all) vertices of X .

This implies that every clique that contains v , v' and some vertex of X is contained in B . Hence, B is the only assembly containing v and v' .

(b) Part (b) follows from (a). \square

LEMMA 2.13. Suppose that the lines L_1 and L_2 meet in a vertex v . Then either every vertex of L_1 has a neighbour on $L_2 \setminus \{v\}$, or every vertex of L_2 has a neighbour on $L_1 \setminus \{v\}$.

PROOF. Assume on the contrary that there exist vertices $v_i \in L_i$ that have no neighbour on $L_{3-i} \setminus \{v\}$, $i = 1, 2$. Then v_1 and v_2 have distance two and thus $c_2 = q^2 + q$ common neighbours. Each common neighbour is the meet of a line of $[v_1, v_2]$ and a line of $[v_2, v_1]$. As $L_1 \in [v_1, v_2]$, $L_2 \in [v_2, v_1]$, and $|[v_1, v_2]| = |[v_2, v_1]| = q + 1$, it follows that v_1 and v_2 have at most $q^2 + 1$ common neighbours. But $c_2 = q^2 + q > q^2 + 1$, a contradiction. \square

LEMMA 2.14.

- (a) A line and an assembly meet in at most $q + 1$ vertices.
- (b) If two lines meet, then their cardinalities differ by at most $q - 1$.
- (c) Every vertex lies on the same number of lines.

PROOF. (a) Consider an assembly B that meets the line L . Then B contains a vertex v with $v \notin L$. As v has at most $q + 1$ neighbours on L , we have $|B \cap L| \leq q + 1$.

(b) Consider two distinct lines L_1 and L_2 containing the vertex v . In view of Lemma 2.13, we can find adjacent vertices $v_1 \in L_1$ and $v_2 \in L_2$ with $v_1, v_2 \neq v$. Let B be the unique assembly with $v_1, v_2 \in B$. Lemma 2.12 shows that $v \in B$ and that $|L_1| - |L_1 \cap B| = |L_2| - |L_2 \cap B|$. As $2 \leq |L_i \cap B| \leq q + 1$, it follows that $|L_1|$ and $|L_2|$ differ by at most $q - 1$.

(c) It suffices to prove this for adjacent vertices v_1 and v_2 . Let s_i be the number of lines of v_i , put $L := v_1v_2$, and denote by l_{ij} , $j = 1, \dots, s_i$, the cardinalities of the lines $\neq L$ on v_i . As v_1 and v_2 have the same number of neighbours, we have

$$\sum_{j=1}^{s_1} (l_{1j} - 1) = \sum_{j=1}^{s_2} (l_{2j} - 1).$$

Assume that $s_2 > s_1$. As every line has at least $\lambda + 2 - (s - 1)q$ vertices, it follows that

$$\sum_{j=1}^{s_1} (l_{1j} - l_{2j}) = \sum_{j=s_1+1}^{s_2} (l_{2j} - 1) \geq (s_2 - s_1)(\lambda + 2 - (s - 1)q).$$

Part (b) shows that $|l_{1j}| \leq |L| + (q - 1)$ and $|l_{2j}| \geq |L| - (q - 1)$. Hence

$$2s_1(q - 1) \geq (s_2 - s_1)(\lambda + 2 - (s - 1)q).$$

As $s_1 \leq s_2 - 1 \leq s - 1$, it follows that

$$2(s - 1)(q - 1) \geq \lambda + 2 - (s - 1)q.$$

Consequently $\lambda + 1 < (3q - 2)(s - 1)$. This contradicts Lemma 2.3. Hence $s_2 \leq s_1$. The same argument shows that $s_1 \leq s_2$. \square

LEMMA 2.15. *Consider two lines L_1 and L_2 on a vertex v . If $|L_1| \geq |L_2|$, then every vertex on $L_1 \setminus \{v\}$ has a neighbour other than v on L_2 .*

PROOF. For $i = 1, 2$, denote by X_i the set consisting of the points of $L_i \setminus \{v\}$ that have a neighbour on $L_{3-i} \setminus \{v\}$. Let B_1, \dots, B_b be the assemblies on v that meet both L_1 and L_2 in a second vertex. Then X_i consists of those points $\neq v$ of L_i that lie in some of the assemblies B_j . Thus

$$|X_i| = \sum_{j=1}^b (|L_i \cap B_j| - 1), \quad i = 1, 2.$$

Put $c := |L_1| - |L_2|$. Lemma 2.12 implies that $c = |L_1 \cap B_j| - |L_2 \cap B_j|$ for every assembly B_j . Hence $|X_1| = |X_2| + bc$.

We have to show that $X_1 = L_1 \setminus \{v\}$, that is $|X_1| = |L_1| - 1$. Assume this does not hold. Then there exists a vertex $v_1 \in L_1$ that has no neighbour $\neq v$ on L_2 . Lemma 2.13 implies that every vertex of L_2 has a neighbour $\neq v$ on L_1 , that is, $|X_2| = |L_2| - 1$. It follows that

$$|L_1| - 1 > |X_1| = |X_2| + bc = |L_2| - 1 + bc.$$

As $b, c \geq 0$, we obtain $|L_1| > |L_2|$, which in turn implies that $c = |L_1| - |L_2| \geq 1$ and $|L_1| \geq |L_2| + b$. Using Lemma 2.14 (b), we obtain $b \leq q - 1$. As, by Lemma 2.14 (a), an assembly and a line share at most $q + 1$ vertices, it follows that

$$|L_2| - 1 = |X_2| = \sum_{j=1}^b (|L_2 \cap B_j| - 1) \leq bq \leq q(q - 1).$$

As $|L_2| \geq \lambda + 2 - (s - 1)q$ (Lemma 2.5), it follows that $\lambda + 1 \leq (s - 1)q + q(q - 1)$. This contradicts Lemma 2.3. \square

LEMMA 2.16. *Consider a vertex v , two lines L and L' and an assembly B with $v \in L, L', B$.*

- (a) *If $|L \cap B| \geq 2$ and $|L| \geq |L'|$, then $|L' \cap B| = |L \cap B| + |L'| - |L|$.*
- (b) *If $|L \cap B| = q + 1$ and $|L'| > |L|$, then $L' \cap B = \{v\}$.*

PROOF. (a) Consider a vertex $w \neq v$ of $B \cap L$. Then B contains all common neighbours of v and w that do not lie on L . As, by Lemma 2.15, w has a neighbour $\neq v$ on L' , it follows that $|L' \cap B| \geq 2$. Lemma 2.12 shows that $|L| - |L \cap B| = \lambda + 2 - |B| = |L'| - |L' \cap B|$.

(b) Assume that $|L' \cap B| \geq 2$. Part (a) shows that $|L' \cap B| = |L \cap B| + |L'| - |L|$. As $|L'| > |L|$ and $|L \cap B| = q + 1$, we obtain $|L' \cap B| > q + 1$. This contradicts Lemma 2.14 (a). \square

LEMMA 2.17. *Suppose that the vertex v is on $t > r$ lines. Then every line on v meets some assembly on v in $q + 1$ vertices, and every assembly on v has at least $tq - r + 2$ vertices.*

PROOF. Denote by $l + 1$ the minimum length of the lines on v . Lemma 2.14 (b) shows that every line on v has at most $l + q$ vertices. As the t lines on v cover the $r\beta$ neighbours of v , it follows that

$$tl \leq r\beta \leq t(l + q - 1).$$

Put $c := \beta - l$. As $r + 1 \leq t \leq 9r/8$ (Lemma 2.6), it follows that

$$\frac{\beta}{r+1} \leq \frac{(t-r)\beta}{t} \leq c \leq \frac{(t-r)\beta}{t} + q - 1 \leq \frac{1}{9}\beta + (q-1).$$

As $s \geq t \geq r + 1$ (Lemma 2.4 (a)) and $\beta \geq q^2 t$ (Lemma 2.3), we also have

$$c \geq \frac{(t-r)\beta}{t} \geq (t-r)q^2 \geq (t-r)q + q^2 - q > (t-r)q + 2q - 3.$$

Denote by \mathcal{B} the set consisting of all assemblies on v , and consider $B \in \mathcal{B}$. Then B meets some line L on v in at least two vertices. As $|L| \leq l + q = \beta + q - c$, Lemma 2.12 shows that

$$\begin{aligned} |B| &= \lambda + 2 + |L \cap B| - |L| \geq \lambda + 4 - |L| \\ &= \beta + 3 + (r-1)(q-1) - |L| \geq (r-2)(q-1) + 2 + c =: b. \end{aligned}$$

As the assemblies of \mathcal{B} cover each of the $r\beta$ neighbours of v exactly once, it follows that $|\mathcal{B}|(b-1) \leq r\beta$. As $c > (t-r)q + 2q - 3$, we also have $b > tq - r + 1$, so every assembly on v has at least $tq - r + 2$ vertices.

Now consider any line L on v . Lemma 2.14 (a) says that $|L \cap B| \leq q + 1$ for every $B \in \mathcal{B}$. Assume that $|L \cap B| \leq q$ for all $B \in \mathcal{B}$. As every vertex of $L \setminus \{v\}$ occurs in an assembly of \mathcal{B} , it follows that $|\mathcal{B}|(q-1) \geq |L| - 1 \geq l$. Hence

$$l(b-1) \leq |\mathcal{B}|(b-1)(q-1) \leq r\beta(q-1).$$

As $l = \beta - c$, it follows that

$$(\beta - c)[(r-2)(q-1) + c] \leq r\beta(q-1),$$

which can be written in the form

$$c[\beta - (r-2)(q-1) - c] \leq 2\beta(q-1).$$

Using the upper and lower bound for c , we obtain

$$\frac{\beta}{r+1} \left[\frac{8}{9}\beta - (r-1)(q-1) \right] \leq 2\beta(q-1).$$

Hence $\frac{8}{9}\beta \leq (3r+1)(q-1)$. As $r+1 \leq t \leq s$, we have $\frac{8}{9}\beta \leq 3s(q-1)$. This contradicts Lemma 2.3. \square

LEMMA 2.18. *Suppose that the vertex v is on $t > r$ lines. Then not all lines on v have the same length.*

PROOF. Assume that every line on v has the same length $l + 1$. Put $l = \beta - c$. As v has $r\beta$ neighbours, we have $t(\beta - c) = tl = r\beta$. It follows that $c = (t-r)\beta/t$.

In view of Lemma 2.17, there exists a line L_0 and an assembly B_0 on v with $|L_0 \cap B_0| = q + 1$. Part (a) of Lemma 2.16 shows that B_0 meets all lines on v in $q + 1$ vertices. Hence $|B_0| = 1 + tq$.

Now consider any assembly B on v . Lemma 2.16 shows that B meets all lines on v in the same number of vertices. Denote this number by $q+1-w$. Then $|B| = 1+t(q+1-w) = |B_0|-tw$. As B meets L_0 in $q+1-w$ vertices and as B_0 meets L_0 in $q+1$ vertices, Lemma 2.12 shows that $|B_0| = |B| + w$. As $|B| = |B_0| - tw$, it follows that $(t-1)w = 0$. Hence $w = 0$.

We have shown that every assembly on v has size $|B_0| = 1 + tq$ and meets all lines on v in $q+1$ vertices.

Consider any neighbour v' of v and denote by B' the assembly on v and v' . Then every line on v' meets B' in at most $q+1$ vertices. As $|B'| = 1 + tq$ and as v' is on t lines (Lemma 2.14 (c)), it follows that all lines on v' meet B' in $q+1$ vertices. Lemma 2.12 (b) implies that all lines on v' have the same length. Hence all lines on v' have length $l+1$. Using an inductive argument, we can show that all lines on a vertex at distance i from v have size $l+1$. Hence all lines of Γ have size $l+1$.

The argument used for the assemblies on v shows that all assemblies have the same size $1 + tq$ and that a line and an assembly either miss or meet in $q+1$ vertices. Thus, if a vertex w has distance one from a line L , then w has precisely $q+1$ neighbours on L (if $w' \in \Gamma_1(w) \cap L$, and if B' is the assembly on w and w' , then $\Gamma_1(w) \cap L = B' \cap L$).

Now consider vertices v_1 and v_2 at distance two. Let u be a common neighbour of v_1 and v_2 , and put $L = v_2u$. As u is a neighbour of v_1 on L , we see that v_1 has $q+1$ neighbours on L . Hence v_1 lies on $q+1$ lines that contain a vertex on L . As v_2 has a neighbour on each of these lines, it has $q+1$ neighbours on each of these lines. Hence $|\Gamma(v_1) \cap \Gamma_1(v_2)| \geq (q+1)^2$, a contradiction. \square

LEMMA 2.19. *Every vertex is on r lines.*

PROOF. By Lemma 2.14, all vertices lie on the same number t of lines. Assume that $t > r$. Consider a vertex v . Denote by $l+1$ the minimum and by $l+1+u$ the maximum size of the lines on v , and denote by m the number of lines of minimum size $l+1$ on v . In view of the preceding lemma, we have $u \geq 1$. We prove the lemma in five steps.

Step 1. $qm \geq tq - r + 1$.

Consider a line L of size $l+1$ on v . In view of Lemma 2.17, there exists an assembly B on v that meets L in $q+1$ vertices. If H is any line on v , then Lemma 2.16 shows that $B \cap H = \{v\}$ if $|H| > |L|$, and $|B \cap H| = q+1$, if $|H| = |L|$. As v lies on m lines of size $l+1$, it follows that $|B| = 1 + mq$. Lemma 2.17 shows that $|B| \geq tq - r + 2$. This proves Step 1.

Step 2. *Every line v has size $l+1$ or $l+2$. If L is a line of size $l+2$ on v , then $L \cap B = \{v\}$ or $|L \cap B| = q+1$ for every assembly B on v .*

Consider any line L_0 of size $l+1+u$ on v and any assembly B on v with $|B \cap L_0| \geq 2$. Put $|B \cap L_0| = q+1-w$. If L is any line on v , then Lemma 2.16 shows that L meets B in $q+1-w+|L|-|L_0|$ vertices. Thus B meets the m lines of size $l+1$ on v in $q+1-w-u$ and every other line on v in at most $q+1-w$ vertices. As v lies on t lines, it follows that

$$|B| - 1 = \sum_{L \in \mathcal{L}} (|B \cap L| - 1) \leq (t-m)(q-w) + m(q-w-u) = t(q-w) - mu.$$

We have $u \geq 1$ and $w \geq 0$. Assume that $u \geq 2$ or $w \geq 1$. As $t \geq m$, it follows that $|B| - 1 \leq tq - 2m$. Using Lemma 2.17 and Step 1, we obtain

$$q(tq - r + 1) \leq q(|B| - 1) \leq tq^2 - 2mq \leq tq^2 - 2(tq - r + 1).$$

This is equivalent to $2qt \leq (r-1)(q+2)$. As $t > r$ and $q \geq 2$, this is a contradiction. Hence $u = 1$, that is every line on v has at most $l+2$ vertices. Furthermore, $w = 0$, that is $|B \cap L| = q+1$.

From now on, we call the lines of size $l + 2$ *long*, and the lines of size $l + 1$ *short*.

Step 3. Consider an assembly B on v . If B meets all long lines on v only in v , then $|B| = 1 + qm$ and B meets every short line on v in $q + 1$ vertices. Otherwise $|B| = 1 + tq - m$, B meets the long lines on v in $q + 1$ and the short lines on v in q vertices.

First consider the case that B meets all long lines on v only in v . Lemma 2.16 shows that B meets all short lines on v in the same number of vertices. Denote this number by $q + 1 - w$ with $w \geq 0$. Then $|B| = 1 + m(q - w)$. Lemma 2.17 implies that $tq + 1 - r \leq m(q - w)$. As $t > r$ and $t \geq m$, it follows that $w < 1$, that is, $w = 0$. Hence $|B| = 1 + mq$, and B meets all short lines on v in $q + 1$ vertices.

Now consider the case that B meets some long line L on v not only in v . Step 2 shows that $|B \cap L| = q + 1$. Lemma 2.16 shows that B meets every line H on v in $q + 1 + |H| - |L|$ vertices, that is, B meets the long lines on v in $q + 1$ and the short lines on v in q vertices. Hence $|B| = 1 + (t - m)q + m(q - 1)$.

Step 4. $(q + 1)m = tq + 1$ and $r\beta = t(l + 1) - m$.

Consider a short line L on v . By Lemma 2.17, there exists an assembly B_1 on v that meets L in $q + 1$ points. The same lemma shows that there exists an assembly B_2 on v that meets some long line on v in $q + 1$ points. Using Step 3, this gives

$$|B_1 \cap L| = q + 1, \quad |B_2 \cap L| = q, \quad |B_1| = 1 + qm, \quad \text{and} \quad |B_2| = 1 + tq - m.$$

Lemma 2.12 shows that $|B_i| + |L| - |B_i \cap L| = \lambda + 2$. Hence $|B_1| = |B_2| + 1$. It follows that $2 + tq - m = 1 + qm$, which gives $(q + 1)m = tq + 1$. As the $t - m$ long and the m short lines on v cover every neighbour of v once, we have $r\beta = b_0 = (t - m)(l + 1) + ml = t(l + 1) - m$. This proves Step 4.

Lemma 2.14 (c) says that every vertex lies on t lines. Hence, everything we have proved for v holds for all vertices. Furthermore, Step 4 show that the integers m and l do not depend on v . It follows that every line has size $l + 1$ or $l + 2$, and that every vertex lies on m short lines.

Step 5. Suppose that the assembly B meets some long line in more than one vertex, and consider any line X with $B \cap X \neq \emptyset$. Then $|B \cap X| = q + 1$ if X is long, and $|B \cap X| = q$ if X is short.

This is true, because Step 3 holds for every vertex.

Now consider a vertex v and two long lines L_1 and L_2 on v . Choose vertices $v_i \in L_i \setminus \{v\}$ such that $d(v_1, v_2) = 2$. Step 5 shows that the assembly on v_1 and v meets L_2 in $q + 1$ vertices. Hence v_1 has $q + 1$ neighbours w_1, \dots, w_{q+1} on L_2 . Denote the line on v_1 and w_i by H_i . W.l.o.g. $w_1 = v$, that is $H_1 = L_1$. Step 5 shows that the assembly on v_2 and w_i meets H_i in at least q vertices. Furthermore, as $H_1 = L_1$ is long, the assembly on v_2 and w_1 meets H_1 in $q + 1$ vertices. It follows that v_2 has at least $q^2 + q + 1$ neighbours on one of the lines H_1, \dots, H_{q+1} . But $d(v_1, v_2) = 2$, so v_1 and v_2 have only $c_2 = q^2 + q$ common neighbours, a contradiction.

We have shown that every vertex is on precisely r lines. Lemma 2.4 implies that every line has $\beta + 1$ vertices, and Lemma 2.9 shows that if a vertex has distance one from a line, then it has q neighbours on that line. As $c_2 = (q + 1)q$, it follows that if two vertices v and w have distance two, then each of the $q + 1$ lines of $[w, v]$ contains q neighbours of v .

Now the proof can be completed using the techniques of Huang [7]. In his paper, he proves a characterization of bilinear forms graphs under certain conditions. Notice that he needs the restriction $q \geq 4$ in section 3 of [7] to apply a well-known Theorem of Buekenhout [3]. However, it was noticed by Cuypers [4] that the use of Buekenhout's result can be avoided, if

one uses a result of Thas and DeClerck [12], which also holds for $q = 2$ and $q = 3$. Huang also uses the so-called 4-vertex condition. However, he needs it only to show that lines have size $\beta + 1$ and that the incidence structure consisting of the vertices and the lines satisfies the dual of Pasch's Axiom, which we have already shown. Thus the arguments of Huang show that q is a prime power, $\beta = q^n - 1$ for some integer $n \geq d$ and Γ is the bilinear forms graph $H_q(n, d)$. This completes the proof of Proposition 2.2. \square

3. GRAPHS WITH CLASSICAL PARAMETERS

Proposition 2.2 requires an integer s satisfying certain conditions. The following lemma gives an existence criterion.

LEMMA 3.1. *There exists an integer s satisfying the hypotheses of Proposition 2.2, provided that one of the following conditions is satisfied.*

- (a) $\beta \geq (2q^4 + 2q^3 + 2q^2 + q - 1)r/(2q - 1)$.
- (b) $q \geq 3$ and $\beta \geq q^{d+3} - 1$.
- (c) $q = 2$ and $\beta \geq q^{d+4} - 1$.

PROOF. As $\lambda + 1 = \beta + (r - 1)(q - 1)$, conditions (b) and (c) in Proposition 2.2 can be written as

$$\beta > \frac{s + 1}{2(s + 1 - r)}[s(q^2 + q - 1) - 2(r - 1)(q - 1)], \text{ and} \quad (1)$$

$$\beta > s(q^3 + q^2 + 2q - 1) - q^2(q^2 + q + 1) - (r - 1)(q - 1). \quad (2)$$

We shall define a real number $s_0 > r$ and then the integer s in such a way that $s_0 - 1 < s \leq s_0$. Then $(s + 1)/(s + 1 - r) \leq s_0/(s_0 - r)$. Hence, instead of inequalities (1) and (2), it suffices to show that

$$\beta > \frac{s_0}{2(s_0 - r)}[s_0(q^2 + q - 1) - 2(r - 1)(q - 1)], \text{ and} \quad (3)$$

$$\beta > s_0(q^3 + q^2 + 2q - 1) - q^2(q^2 + q + 1) - (r - 1)(q - 1). \quad (4)$$

We choose $s_0 := 2qr/(2q - 1)$. Then $s_0/(s_0 - r) = 2q$. Hence, (3) and (4) are satisfied provided that

$$\beta \geq s_0(q^3 + q^2 + 2q - 1) - r(q - 1) = \frac{(2q^4 + 2q^3 + 2q^2 + q - 1)r}{2q - 1}.$$

This proves part (a). Part (b) is an immediate consequence of part (a). In order to prove part (c), we have to choose s_0 much more carefully. First notice that in the case $q = 2$ and $d = 3$, we can choose $s = r = 7$; then (1) and (2) and hence hypotheses (a), (b) and (c) of Proposition 2.2 are satisfied. Suppose from now on that $q = 2$ and $d \geq 4$. Let u be the larger root of the equation

$$x(5x - 2) = 2(x - 1)(15x - 1).$$

Then $u \approx 1.129$ and

$$c := \frac{x(5x - 2)}{2(x - 1)} = (15x - 1) \approx 15.94.$$

We choose $s_0 := ur$. As $q = 2$, inequalities (3) and (4) require that

$$\beta > \frac{u}{2(u - 1)}(5ur - 2r + 2) = cr + \frac{u}{u - 1}$$

and

$$\beta > 15ur - 27 - r = cr - 27.$$

If $\beta \geq q^{d+4} - 1$, then $\beta > 16(q^d - 1) + 15 = 16r + 15$ and (3) and (4) are satisfied. This proves part (c). \square

Theorem 1.1 is an immediate consequence of Proposition 2.2 and part (b) and (c) of the preceding lemma. It remains to prove Corollary 1.3.

PROOF OF COROLLARY 1.3. Suppose that Γ is a distance-regular graph of diameter $d \geq 3$. Then q is an integer with $q \neq 0, -1$ (see Theorem 6.2.1 in [2]). Put $r := \begin{bmatrix} d \\ 1 \end{bmatrix}$.

First consider the case that $q \leq -2$. Then $\alpha < 0$, as $c_2 = (1 + q)(1 + \alpha)$. If d is even, then $r = (q^d - 1)/(q - 1) < 0$. As $0 < b_0 = r\beta$ and $0 \leq \lambda = \beta - 1 + (r - 1)\alpha$, it follows $1 - \alpha(r - 1) \leq \beta < 0$. If d is odd, then $r > 0$ and hence $\beta > 0$. In this case $0 \leq a_2 = b_0 - b_2 - c_2 = (q + 1)(\beta + r\alpha - 2\alpha - q\alpha - 1)$ gives $0 \leq \beta \leq 1 - \alpha(r - q - 2)$. Hence, for d even or odd and $q \leq -2$, we have $|\beta| \leq 1 + |\alpha|(|r| + |q| - 1)$.

Now consider the case when $q \geq 1$ and $\alpha \neq q - 1, q$. Theorem 8.4.3 in [2] says that the multiplicity of the second largest eigenvalue of Γ is given by

$$f_1 = \frac{qr\beta \left(1 + \begin{bmatrix} d-2 \\ 1 \end{bmatrix} \alpha + q^{d-2}\beta\right)}{(\beta - \alpha + q) \left(1 + \alpha \begin{bmatrix} d-1 \\ 1 \end{bmatrix}\right)}.$$

As $c_2 = (1 + q)(\alpha + 1)$ and $a_1 = \beta - 1 + \alpha(r - 1)$ are integers, we see that $\alpha' := \alpha(q + 1)$ and $\beta' := \beta(q + 1)$ are integers. This implies that $\beta' - \alpha' + q(q + 1)$ is an integer that divides the integer $qr\beta'(q + 1 + \begin{bmatrix} d-2 \\ 1 \end{bmatrix} \alpha' + q^{d-2}\beta')$. This implies that $\beta' - \alpha' + q(q + 1)$ divides

$$\begin{aligned} & qr(\alpha' - q(q + 1)) \left(q + 1 + \begin{bmatrix} d-2 \\ 1 \end{bmatrix} \alpha' + q^{d-2}(\alpha' - q(q + 1)) \right) \\ &= (q + 1)^2 r(r - 1)(\alpha - q)(\alpha + 1 - q). \end{aligned}$$

If $\alpha \neq q - 1, q$, this yields the bound of Corollary 1.3, because $\beta' - \alpha' + q(q + 1) = (q + 1)(\beta - \alpha + q)$.

Next consider the case $q \geq 1$ and $\alpha = q - 1$. If $q = 1$, and $\beta > 3$, then it was shown by Egawa [5] that Γ is a Hamming graph. If $q \geq 2$, then part (a) Lemma 3.1 shows that Γ is a bilinear forms graph or that $\beta \leq (2q^4 + 2q^3 + 2q^2 + q - 1)r/(2q - 1)$.

Finally, consider the case that $q \geq 1$ and $\alpha = q$. If $q = 1$, then it was shown independently by Neumaier [10] and Terwilliger [11] that Γ is a Johnson graph. If $q \geq 2$ and $\beta > 8r(q^2 + 2q)/3$, then it follows easily from Theorem 2.3 in [9] that Γ is a Grassmann graph (in Theorem 2.3 of [9], define s to be the unique integer satisfying $4r/3 - 1 < s \leq 4r/3$). \square

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